(19) Europäisches Patentamt
European Patent Office
Office européen des brevets

(11) Publication number: **0 603 916 A2**

(12) EUROPEAN PATENT APPLICATION

(21) Application number: **93120960.5**

(51) Int. Cl.5: **H04L 12/56**, H04Q 11/04

(22) Date of filing: **27.12.93**

(30) Priority: **25.12.92 JP 345020/92**

(43) Date of publication of application:
29.06.94 Bulletin 94/26

(84) Designated Contracting States:
DE FR GB IT SE

(71) Applicant: **NEC CORPORATION**
7-1, Shiba 5-chome
Minato-ku
Tokyo 108-01(JP)

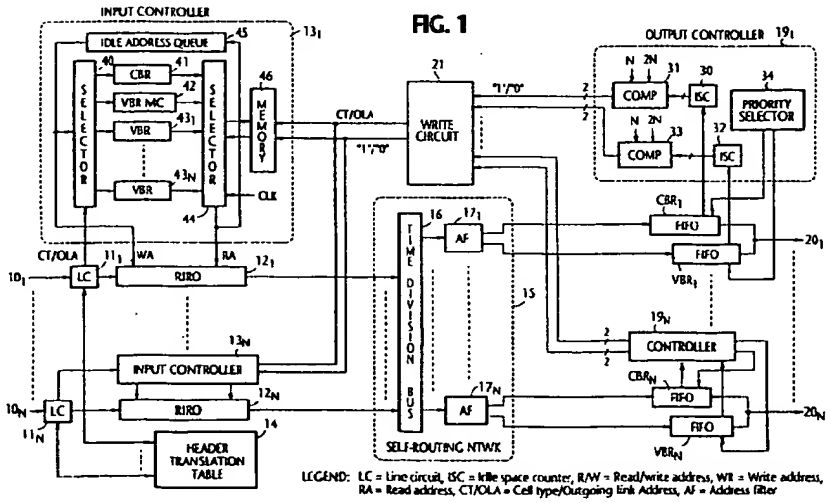
(72) Inventor: **Fan, Ruixue, c/o NEC CORPORATION**
7-1, Shiba 5-chome
Minato-ku, Tokyo(JP)

(74) Representative: **VOSSIUS & PARTNER**
Postfach 86 07 67
D-81634 München (DE)

(54) Packet switching system using idle/busy status of output buffers.

(57) A packet switching system comprises input buffers and output buffers, and a self-routing network for routing packets from the input buffers to the output buffers according to the routing information contained in their header. An input controller (13) associated with each input buffer, includes an idle/busy memory. An output controller (19), associated with each output buffer, detects the idle space of the associated output buffer and determines if the idle space is greater than a predetermined value. If the idle space is determined to be greater than the predetermined value, an idle status bit is generated and if the idle space is determined to be smaller than the predetermined value, a busy status bit is generated. The idle and busy status bits are broadcast to the idle/busy memory of all input controllers. Each input controller (13) writes an incoming packet into the associated input buffer and reads it onto the self-routing network if an idle status bit is stored in the idle/busy memory corresponding to the routing information of the incoming packet.

FIG. 1



EP 0 603 916 A2

BNSDOCID: <EP 0603916A2 I >

The present invention relates generally to packet switching systems, and more specifically to a packet switching system of the input-output buffering type for switching packets such as asynchronous transfer mode (ATM) cells.

In a packet switching system of the input and output buffering type, as described in a paper "A very high-speed ATM switch with input and output buffers", Yukihiko Do et al, ISS'92 pages 231-235, incoming packets are stored In an input buffer and launched onto a self-routing network where it is routed to a destination output buffer according to thc routing information of the packets. To achieve high throughput, it is necessary to provide as many high capacity output buffers as possible for bursty traffic and to provide a high speed scheduling algorithm for multicast traffic. It is still desirable to reduce the amount of output buffers required and to simplify the scheduling algorithm for both point-to-point and multicast traffic.

It is therefore an object of the present invention to provide a high throughput packet switching system which eliminates the need for high capacity output buffers.

Another object of the present invention is to provide a high throughput packet switching system which can be implemented with a simplified scheduling algorithm for both point-to-point and multicast traffic.

According to a broader aspect of the present invention, there is provided a packet switching system comprising a plurality of input buffers connected respectively to a plurality of incoming links and a plurality output buffers connected respectively to a plurality of outgoing links. A self-routing network is connected between the input and output buffers for routing packets from the input buffers to thc output buffers according to routing information contained in their header. A plurality of output controllers are associated respectively with the output buffers. Each output controller determines the idle space of the associated output buffer which is available for storing additional packets and determines whether the idle space is greater than a predetermined value. If the idle space is determined to be greater than the predetermined value, an idle status bit is generated and if the idle space is determined to be smaller than the predetermined value, a busy status bit is generated. A plurality of input controllers are associated respectively with the input buffers for storing the idle and busy status bits from the output controllers in an idle/busy memory and writing an incoming packet into the associated input buffer. Each input controller reads a packet from the associated input buffer onto the self-routing network if an idle status bit is stored in the idle/busy memory corresponding to the destination of the stored packet.

According to a second aspect of the present invention, there is provided a packet switching system which comprises input buffers connected respectively to incoming links, output buffers connected respectively to outgoing links, and a self-routing network for routing packets from the input buffers to the output buffers according to routing information contained in the packets. A plurality of output controllers are associated respectively with the output buffers for determining an idle space of the associated output buffer which is available for storing additional point-to-point packets and additional multicast packets. Each output controller determines whether the idle space of point-to-point packets is greater than a first predetermined value and determines whether the idle space is greater than a second predetermined value smaller than the first predetermined value, and generates a first idle status bit or a first busy status bit into an idle/busy status memory if the idle space is determined to be greater or smaller than the first predetermined value, respectively, and generates a second idle status bit or a second busy status bit if the idle space is determined to be greater or smaller than the second predetermined value, respectively. A plurality of input controllers are associated respectively with the input buffers for storing the first idle and busy status bits and the second idle and busy status bits in an idle/busy memory and storing an incoming point-to-point packet into the associated input buffer and reading it if a first idle status bit is stored in the idle/busy memory corresponding to the destination of the point-to-point packet, and writing an incoming multicast packet into the associated input buffer and reading it if a second idle status bit is stored in the idle/busy memory corresponding to the incoming multicast packet.

BRIEF DESCRIPTION OF THE DRAWINGS

The present invention will be described in further detail with reference to the accompanying drawings, in which:

Fig. 1 is a block diagram of an ATM self-routing switching system according to the present invention;

Fig. 2 is a table illustrating the idle/busy status bits generated by an output controller according to different cell types and stored in the idle/busy status memory of each input controller;

Fig. 3 is a flowchart of a sequence of instructions performed by a selector of an input controller when a cell is read out of an associated RIRO memory; and

Fig. 4 is a modified version of the flowchart of Fig. 3.

BNSDOCID: <EP 0603916A2 I >

DETAILED DESCRIPTION

Referring now to Fig. 1, there is shown an asynchronous transfer mode (ATM) switching system according to the present invention. The switching system comprises a plurality of line circuits 11_1 ~ 11_N which respectively serve as interfaces between incoming links 10_1 ~ 10_N and input buffers, or random-in-random-out memories 12_1 ~ 12_N. Connected respectively to the line circuits 11_1 ~ 11_N and RIRO memories 12_1 ~ 12_N are input controllers 13_1 ~ 13_N, each of which provides address control on incoming ATM cells to determine their write and read addresses of the associated RIRO memory in a manner as will be described in detail later. A header translation table **14** is connected to the line circuits 11_1 ~ **11N** to translate the VCI/VPI (virtual call identifier/virtual path identifier) field contained in the header of each incoming packet, called "ATM cell" to an outgoing link address (OLA) that specifies the appropriate outgoing link of the switching system in a well known manner. Each incoming cell contains a cell type field in its header to indicate the type of the cell as constant bit rate (CBR), variable bit rate (VBR), a point-to-point, or multicast data or a combination of these. Following the header translation on an incoming ATM cell, each line circuit 11_i (where i = 1, 2, ... N) supplies the cell type (CT) data and the outgoing link address (OLA) to the associated input controller 13_i and the whole contents of the cell to the associated RIRO memory 12_i where the cell is stored in a location specified by a write address WA supplied from input controller 13_i and eventually read out of the RIRO memory 12_i in response to a read address RA from input controller 13_i.

The outputs of RIRO memories 12_1 ~ 12_N are connected to a self-routing network **15** which includes a time division bus **16**, on which the outputs of RIRO memories **12** are time-division multiplexed, and a plurality of address filters 17_1 ~ 17_N connected to the bus **16**. Address filters 17_i are provided corresponding respectively to the outgoing links 20_i and are assigned the output link address (OLA) of the corresponding output link plus two cell type identifiers identifying respectively CBR and VBR cells. Each address filter **17** has two output terminals for respectively delivering CBR and VBR cells to the output buffers of a corresponding pair. Each address filter allows passage of those ATM cells to one of the output buffers of the corresponding pair if it contains the output link address of the address filter and one of the cell type identifiers of the address filter. The outgoing buffers are implemented with first-in-first-out memories $\mathbf{CBR}_1$ ~ $\mathbf{CBR}_N$ for exclusively storing CBR cells and FIFO memories $\mathbf{VBR}_1$ ~ $\mathbf{VBR}_N$ for exclusively storing VRB cells. The outputs of the FIFO memories of each pair are connected together to a corresponding one of outgoing link 20_1 ~ 20_N.

A plurality of output controllers 19_1 ~ 19_N are connected to outgoing FIFO memories $\mathbf{CBR}_1$ ~ $\mathbf{CBR}_N$ and $\mathbf{VBR}_1$ ~ $\mathbf{VBR}_N$, respectively. Each output controller includes an idle space (cell) counters **30** and **32** respectively connected to the associated FIFO memories to identify the type of cells stored in the associated FIFO memory using their cell type field and detect the number of cells of particular type that can be stored in each of the associated FIFO memories. Comparators **31** and **33** are connected to the output of idle space counters **30** and **32**, respectively, for making a comparison between the idle space count from the respective idle space counter with reference values N and 2N to produce a logic-0 (idle) which indicates that a destination FIFO memory has a sufficient amount of idle space for accepting cells or a logic logic-1 indicating that it has not.

More specifically, comparator **31** produces a 0 if the idle space count in the FIFO memories $\mathbf{CBR}_i$ (where i = 1, 2, ..., N) is equal to or greater than N or a 1 if the idle space count is smaller than N. Comparator **31** further produces a 0 if the idle space count in the FIFO memories $\mathbf{CBR}_i$ is equal to or greater than 2N or a 1 if the idle space count is smaller than 2N. Likewise, comparator **33** produces a 0 if the idle space count in the FIFO memories $\mathbf{VBR}_i$ is equal to or greater than N or a 1 if the idle space count is smaller than N. Comparator **33** further produces a 0 if the idle space count in the FIFO memories $\mathbf{VBR}_i$ is equal to or greater than 2N or a 1 if the idle space count is smaller than 2N.

A priority selector **34** is further provided in each of the output controllers **19**. Priority selector **34** is connected to the FIFO memories of the associated pair for reading a CBR cell first from the CBR memory of the pair if a VBR cell is simultaneously stored in the VBR memory of the pair.

The comparators **31** and **33** of each output controller supply their idle/busy status bits to a write circuit **21**. Write circuit **21** sequentially selects the output controllers to receive their output signals and produces a write address signal and broadcasts it together with the idle/busy status bits to all input controllers **13**.

Each input controller **13** includes an input selector **40** and an output selector **44**. Connected between the input and output selectors **40** and **44** are a plurality of FIFO memories **41**, **42**, 43_1 ~ 43_N for respectively storing different types (priority and destination) of cells. An idle address queue **45** is connected between an output terminal of selector **44** and an Input of selector **40** for storing idle

BNSDOCID: <EP 0603916A2 I >

addresses that can be used as a write address for each incoming cell. Further provided in each input controller is an idle/busy memory **46** for storing the idle/busy status bits broadcast from the write circuit **21** as shown In Fig. 2.

The selector **40** of each input controller receives the CT and OLA fields of an incoming cell from the associated line circuit **11**. On receiving the CT/OLA data, selector **40** reads an idle address from queue **45** and supplies it to the write address input of the associated RIRO memory **12** for writing the incoming cell into the memory **12**. At the same time, it appends the CT/OLA data to the write address to form a read address signal and supplies it to one of the memories **41**, **42**, 43_1 ~ 43_N according to the type of the cell now stored in the RIRO memory **12**.

CBR cells of either point-to-point or multicast type are stored in queue **41**. VBR multicast cells are stored in queue **42** and VBR point-to-point cells are stored into one of the memories 43_1 ~ 43_N according to their destinations indicated by their OLA data. Note that CBR cells are given highest priority, VBR multicast cells second priority, and VBR point-to-point cells lowest priority.

The operation of the output selector **44** will be described with reference to Fig. 3. Selector **44** receives the system clock to determine the slot timing for reading an address signal (step **50**), set a variable Q to zero (step **51**), and accesses the output end of queues **41**, **42** and **43**, starting with queue **41** to determine whether an address signal is stored in that queue (step **52**). If there is one, control branches at step **52** to step **53** to read the stored address signal and apply the CT and OLA data of the address signal to the idle/busy memory **46**. An idle/busy status bit is therefore read out of a location of the idle/busy memory **46** corresponding to the cell to be read out of the associated RIRO memory **12**. At step **54**, the variable Q is incremented by 1 and, at step **55**, control determines whether the idle/busy status bit from the idle/busy memory **46** is 1 or 0. If the idle/busy status bit is 1, control returns to step **52** to repeat the process on the output end of the next queue, and if it is 0, control branches to step **56** to supply the address data retrieved at step **53** to the associated RIRO memory **12** as a read address and to idle address queue **45** as an idle address. Control proceeds to step **57** to, check to see if the variable Q is equal to K which indicates the number of memories **41**, **42** and **43**. If Q < K, control returns to step **52** to repeat the above process on the next queue, and if Q = K, it returns to step **50**. If the decision at step **52** is negative, step **58** is executed by incrementing the variable Q by 1 and the value of Q is then checked at step **56**.

In this way, those cells on which decision is made that their destination output buffers are not sufficient to accept are left in their memories until the destination output buffers are made available.

The advantageous effect of the present invention is that the storage capacity of each output buffer can be reduced and scheduling of cells to be launched into the self-routing network **15** from each input buffer can be performed at relatively low speed.

A modified embodiment is shown in Fig. 4 in which priorities are assigned to memories **43** by taking into account the number of address signals stored in each of these memories **43** as well as the cell type of the stored address signals. As illustrated, step **60** is provided between steps **51** and **52** of flowchart of Fig. 3 for counting the address signals stored In each of the memories **41**, **42**, 43_1 ~ 43_N and priorities are respectively assigned to these memories according to their stored address counts and their cell types, so that address signals are read from these memories according to the assigned priorities at step **52**.

Claims

1. A packet switching system comprising:

a plurality of input buffers connected respectively to a plurality of incoming links;

a plurality output buffers connected respectively to a plurality of outgoing links;

a routing network connected between said input buffers and said output buffers for routing packets from the input buffers to the output buffers according to routing information contained in the packets;

a plurality of output controllers associated respectively with said output buffers, each of the output controllers determining an idle space of the associated output buffer which is available for storing additional packets and determining whether the idle space is greater than a predetermined value, and generating an idle status bit if the idle space is determined to be greater than said predetermined value or a busy status bit if the idle space is determined to be smaller than said predetermined value; and

a plurality of input controllers associated respectively with said input buffers, each of the input controllers including memory means for storing said idle status bit and said busy status bit from said output controllers, and means for writing an incoming packet into the associated input buffer and reading the incoming packet therefrom onto the routing network if an idle status bit is stored in said memory means corresponding to the routing information of

BNSDOCID: <EP 0603916A2 I >

said incoming packet.

2. A packet switching system comprising:

N input buffers connected respectively to N incoming links;

N output buffers connected respectively to N outgoing links;

a self-routing network connected between said input buffers and said output buffers for routing packets from the input buffers to the output buffers according to routing information contained in the packets;

N output controllers associated respectively with said output buffers, each of the output controllers determining an idle space of the associated output buffer which is available for storing additional packets and determining whether the idle space is greater than N, and generating an idle status bit if the idle space is determined to be greater than N or a busy status bit if the idle space is determined to be smaller than N; and

N input controllers associated respectively with said input buffers, each of the input controllers including memory means for storing said idle status bit and said busy status bit from said output controllers, and means for writing an incoming packet into the associated input buffer and reading the incoming packet therefrom onto the self-routing network if an idle status bit is stored in said memory means corresponding to the routing information of said incoming packet.

3. A packet switching system comprising:

a plurality of input buffers connected respectively to a plurality of incoming links;

a plurality of output buffers connected respectively to a plurality of outgoing links;

a self-routing network connected between said input buffers and said output buffers for routing packets from the input buffers to the output buffers according to routing information contained in the packets;

a plurality of output controllers associated respectively with said output buffers, each of the output controllers determining an idle space of the associated output buffer which is available for storing additional point-to-point packets and additional multicast packets, determining whether the idle space of point-to-point packets is greater than a first predetermined value, determining whether the idle space of multicast packets is greater than a second predetermined value smaller than said first predetermined value, generating a first idle status bit or a first busy status bit if the idle space is determined to be greater or smaller than said first predetermined value, respectively, and generating a second idle status bit or a second busy status bit if the idle space is determined to be greater or smaller than said second predetermined value, respectively; and

a plurality of input controllers associated respectively with said input buffers, each of the input controllers including memory means for storing said first idle and busy status bits and second idle and busy status bits from said output controllers, and means for writing an incoming point-to-point packet into the associated input buffer and reading the incoming point-to-point packet therefrom onto the self-routing network if a first idle bit from is stored in said memory means corresponding to the routing information of the incoming point-to-point packet, storing an incoming multicast packet into the associated input buffer and reading the incoming multicast packet therefrom onto the self-routing network if a second idle status bit is stored in said memory means corresponding to the incoming multicast packet.

4. A packet switching system comprising:

N input buffers connected respectively to N incoming links;

N output buffers connected respectively to N outgoing links;

a self-routing network connected between said input buffers and said output buffers for routing packets from the input buffers to the output buffers according to routing information contained in the packets;

N output controllers associated respectively with said output buffers, each of the output controllers determining an idle space of the associated output buffer which is available for storing additional point-to-point packets and additional multicast packets, determining whether the idle space of point-to-point packets is greater than M greater than N and whether the idle space of multicast packets is greater than N, generating a first idle status bit or a first busy status bit if the idle space is determined to be greater or smaller than M, respectively, and generating a second idle status bit or a second busy status bit if the idle space is determined to be greater or smaller than N, respectively; and

N input controllers associated respectively with said input buffers, each of the input controllers including memory means for storing the first idle and busy status bits and the second idle and busy status bits from said output controllers, and means for writing an incoming point-to-point packet into the asso-

BNSDOCID: <EP 0603916A2 I >

ciated input buffer and reading the incoming point-to-point packet therefrom onto the self-routing network if a first idle status bit is stored in said memory means corresponding to the routing information of the incoming point-to-point packet, writing an incoming multicast packet into the associated input buffer and reading the incoming multicast packet therefrom onto the self-routing network if a second idle status bit is stored in said memory means corresponding to the incoming multicast packet.

5. A packet switching system as claimed in claim 1, 2, 3 or 4, wherein each of said input controllers comprises:

an idle address queue;

a high priority queue and a low priority queue;

first selector means for retrieving an address signal from said idle address queue and writing an incoming packet into a location of the associated input buffer specified by the retrieved address signal, and writing the retrieved address signal into one of said high and low priority queues according to priority of the incoming packet;

second selector means for retrieving an address signal from each of said high and low priority queues, retrieving an idle/busy status bit from said memory means corresponding to the retrieved address signal, determining whether the retrieved idle/busy status is 0 or 1, and supplying the retrieved address signal to said idle address queue and to the associated input buffer when the retrieved idle/busy status bit is determined to be 0 for reading a packet from the associated input buffer onto the self-routing network, or leaving the retrieved address signal in said first and second priority queues if the retrieved idle/busy status bit is determined to be 1,

each of said output buffers including a high priority memory for storing packets of high priority and a low priority memory for storing packets of low priority,

each of the output controllers including means for forwarding a packet from the high priority memory of the associated output buffer to the associated outgoing link earlier than a packet from the low priority memory of the associated output buffer when more than one packet exists simultaneously in said high and low priority memories of the associated output buffer.

6. A packet switching system as claimed in claim 1, 2, 3 or 4, wherein each of said input controllers comprises:

an idle address queue;

a plurality of busy address queues;

first selector means for retrieving an address signal from said idle address queue and writing an incoming packet into the associated input buffer corresponding to the retrieved address signal, and writing the retrieved address signal into one of said busy address queues according to the routing information of the incoming packet; and

second selector means for retrieving an address signal from each of said busy address queues, retrieving an idle/busy status bit from said memory means corresponding to the address signal retrieved from said busy address queues, determining whether the retrieved idle/busy status is 0 or 1, and supplying the retrieved address signal to said idle address queue and to the associated input buffer If the retrieved idle/busy status bit is determined to be 0 for reading a packet from the associated input buffer onto the self-routing network, or leaving the retrieved address signal in said busy address queues if the retrieved idle/busy status bit is determined to be 1.

7. A packet switching system as claimed in claim 1, 2, 3 or 4, wherein each of said input controllers comprises:

an idle address queue;

a plurality of busy address queues;

first selector means for retrieving an address signal from said idle address queue and writing an incoming packet into the associated input buffer corresponding to the retrieved address signal, and writing the retrieved address signal into one of said busy address queues according to the routing information of the incoming packet; and

second selector means for retrieving an address signal from said busy address queues according priority of packets stored therein, retrieving an idle/busy status bit from said memory means corresponding to the address signal retrieved from said busy address queues, determining whether the retrieved idle/busy status is 0 or 1, and supplying the retrieved address signal to said idle address queue and to the associated input buffer if the retrieved idle/busy status bit is determined to be 0 for reading a packet from the associated input buffer onto the self-routing network, or leaving the retrieved address signal in said busy address queues if the retrieved idle/busy status bit is determined to be 1.

BNSDOCID: <EP 0603916A2 I >

8. A packet switching system as claimed In claim 1, 2, 3 or 4, further comprising means for sequentially receiving an idle/busy status bit from each of said output controllers and generating a write address signal and broadcasting the write address signal together with the idle/busy status to all of said input controllers for writing the idle/busy status bit into the memory means of said input controllers.

BNSDOCID: <EP 0603916A2 I >

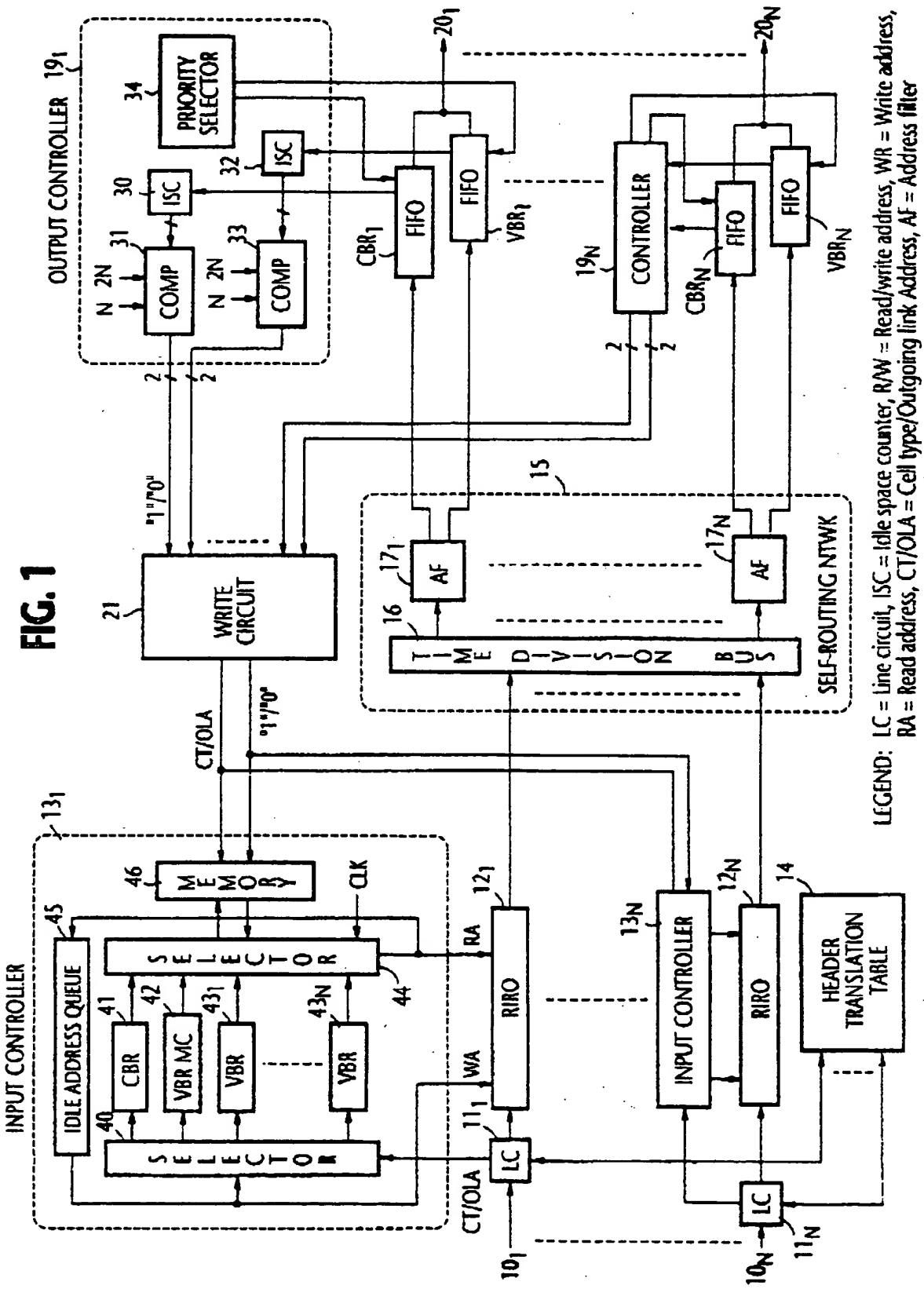

OUTPUT CONTROLLER 19₁
PRIORITY SELECTOR
34
ISC
32
ISC
30
COMP
31
N 2N
COMP
33
N 2N
2
2
"1"/"0"
FIFO
CBR₁
FIFO
VBR₁
20₁
20N
19N CONTROLLER
FIFO
CBRN
FIFO
VBRN
2
2
15
SELF-ROUTING NTWK
TIME DIVISION BUS
16
17₁ AF
17N AF
21
WRITE CIRCUIT
"1"/"0"
CT/OLA
13₁
INPUT CONTROLLER
IDLE ADDRESS QUEUE
45
MEMORY
46
CLK
SELECTOR
44
40
SELECTOR
CBR
41
VBR MC
42
VBR
43₁
VBR
43N
RA
WA
12₁
RIRO
11₁
LC
CT/OLA
10₁
13N
INPUT CONTROLLER
12N
RIRO
14
HEADER TRANSLATION TABLE
LC
11N
10N
LEGEND: LC = Line circuit, ISC = Idle space counter, R/W = Read/write address, WR = Write address, RA = Read address, CT/OLA = Cell type/Outgoing link Address, AF = Address filter

FIG. 1

FIG. 2

	CELL TRANSFER MODE	MULTICAST			POINT-TO-POINT			MULTICAST			POINT-TO-POINT		
	FIFO	CBR_1	----	CBR_N	CBR_1	-----	CBR_N	VBR_1	-----	VBR_N	VBR_1	-----	VBR_N
OUTPUT OF COMP-PARATOR 31	IDLE SPACE COUNT ≥ N	0	----	0		----			----			----	
	IDLE SPACE COUNT < N	1	----	1		----			----			----	
	IDLE SPACE COUNT ≥ 2N		----		0	----	0		----			----	
	IDLE SPACE COUNT < 2N		----		1	----	1		----			----	
OUTPUT OF COMP-PARATOR 33	IDLE SPACE COUNT ≥ N		----			----		0	----	0		----	
	IDLE SPACE COUNT < N		----			----		1	----	1		----	
	IDLE SPACE COUNT ≥ 2N		----			----			----		0	----	0
	IDLE SPACE COUNT < 2N		----			----			----		1	----	1

LEGEND: 0 = IDLE (FIFO HAS SUFFICIENT IDLE SPACE TO ACCEPT INCOMING CELLS)
1 = BUSY (FIFO HAS NOT SUFFICIENT IDLE SPACE TO ACCEPT INCOMING CELLS)

EPÜ: <EP 0603916A2 I >

FIG. 3

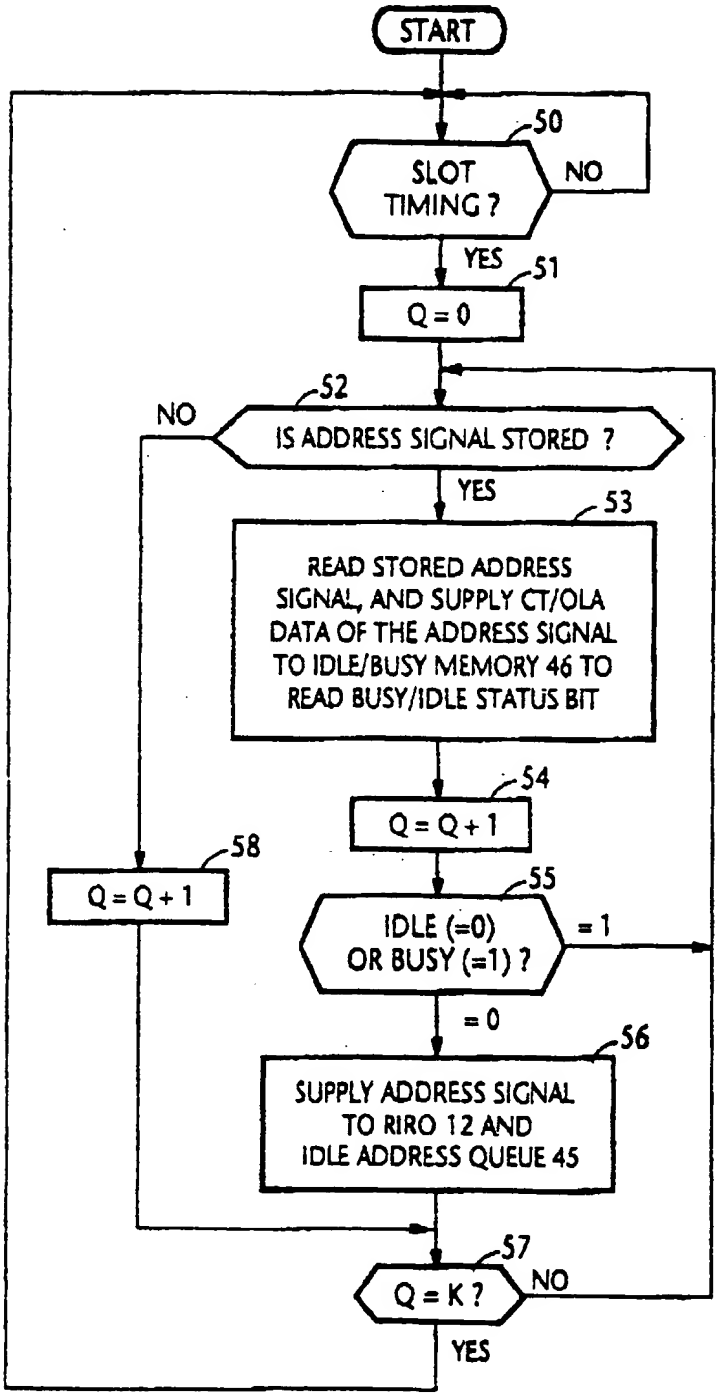
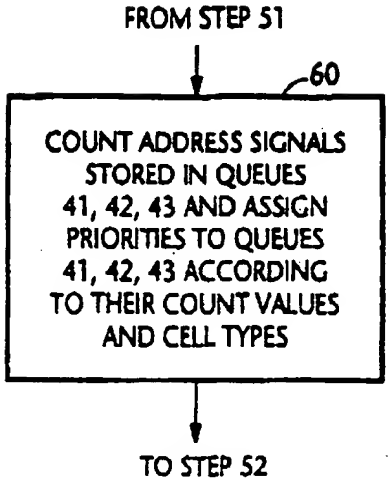

START
50
SLOT TIMING ?
NO
YES
51
Q = 0
52
NO
IS ADDRESS SIGNAL STORED ?
YES
53
READ STORED ADDRESS SIGNAL, AND SUPPLY CT/OLA DATA OF THE ADDRESS SIGNAL TO IDLE/BUSY MEMORY 46 TO READ BUSY/IDLE STATUS BIT
54
Q = Q + 1
58
Q = Q + 1
55
IDLE (=0) OR BUSY (=1) ?
= 1
= 0
56
SUPPLY ADDRESS SIGNAL TO RIRO 12 AND IDLE ADDRESS QUEUE 45
57
Q = K ?
NO
YES

FIG. 4


FROM STEP 51
60
COUNT ADDRESS SIGNALS STORED IN QUEUES 41, 42, 43 AND ASSIGN PRIORITIES TO QUEUES 41, 42, 43 ACCORDING TO THEIR COUNT VALUES AND CELL TYPES
TO STEP 52

(19) Europäisches Patentamt
European Patent Office
Office européen des brevets

(11) EP 0 603 916 A3

(12) EUROPEAN PATENT APPLICATION

(88) Date of publication A3:
27.12.1996 Bulletin 1996/52

(43) Date of publication A2:
29.06.1994 Bulletin 1994/26

(21) Application number: **93120960.5**

(22) Date of filing: **27.12.1993**

(51) Int. Cl.[6]: **H04L 12/56**, H04Q 11/04

(84) Designated Contracting States:
DE FR GB IT SE

(30) Priority: **25.12.1992 JP 345020/92**

(71) Applicant: **NEC CORPORATION**
Tokyo (JP)

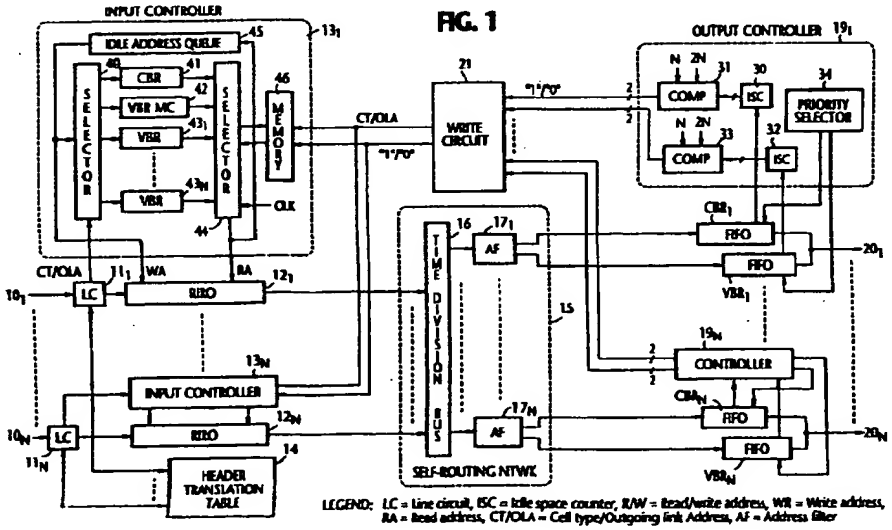
(72) Inventor: **Fan, Ruixue,**
c/o NEC CORPORATION
Minato-ku, Tokyo (JP)

(74) Representative: **VOSSIUS & PARTNER**
Postfach 86 07 67
81634 München (DE)

(54) Packet switching system using idle/busy status of output buffers

(57) A packet switching system comprises input buffers and output buffers, and a self-routing network for routing packets from the input buffers to the output buffers according to the routing information contained in their header. An input controller (13) associated with each input buffer, includes an idle/busy memory. An output controller (19), associated with each output buffer, detects the idle space of the associated output buffer and determines if the idle space is greater than a predetermined value. If the idle space is determined to be greater than the predetermined value, an idle status bit is generated and if the idle space is determined to be smaller than the predetermined value, a busy status bit is generated. The idle and busy status bits are broadcast to the idle/busy memory of all input controllers. Each input controller (13) writes an incoming packet into the associated input buffer and reads it onto the self-routing network if an idle status bit is stored in the idle/busy memory corresponding to the routing information of the incoming packet.

FIG. 1



EP 0 603 916 A3

European Patent Office

EUROPEAN SEARCH REPORT

Application Number
EP 93 12 0960

DOCUMENTS CONSIDERED TO BE RELEVANT

Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int.Cl.5)
A	IEEE GLOBAL TELECOMMUNICATIONS CONFERENCE - GLOBECOM '91, vol. 1 OF 3, 2 - 5 December 1991, PHOENIX, ARIZONA, pages 347-351, XP000325996 BADRAN H F ET AL: "HEAD OF LINE ARBITRATION IN ATM SWITCHES WITH INPUT-OUTPUT BUFFERING AND BACKPRESSURE CONTROL" * figure 1 * * abstract * ---	1,2	H04L12/56 H04Q11/04
A	ELECTRONICS & COMMUNICATIONS IN JAPAN, PART I - COMMUNICATIONS, vol. 74, no. 11, 1 November 1991, pages 17-24, XP000305770 RUIXUE FAN ET AL: "AN INPUT BUFFER-TYPE ATM SWITCHING SYSTEM USING SCHEDULE COMPARISON" * figure 1 * * Page 18, chapter 2 - end of page * ---	1,2	TECHNICAL FIELDS SEARCHED (Int.Cl.5)
A	IEEE GLOBAL TELECOMMUNICATIONS CONFERENCE GLOBECOM '91, vol. 3 of 3, 2 - 5 December 1991, PHOENIX, ARIZONA, pages 1611-1618, XP000272782 CHAO H J: "A NOVEL ARCHITECTURE FOR QUEUE MANAGEMENT IN THE ATM NETWORK" * figure 8 * Chapter III, paragraphs A,B,C --- -/--	5,6	H04L H04Q

The present search report has been drawn up for all claims

Place of search	Date of completion of the search	Examiner
THE HAGUE	29 October 1996	Dhondt, E

CATEGORY OF CITED DOCUMENTS

X : particularly relevant if taken alone
Y : particularly relevant if combined with another document of the same category
A : technological background
O : non-written disclosure
P : intermediate document

T : theory or principle underlying the invention
E : earlier patent document, but published on, or after the filing date
D : document cited in the application
L : document cited for other reasons

& : member of the same patent family, corresponding document

EPO FORM 1503 03.82 (P04C01)

European Patent Office

EUROPEAN SEARCH REPORT

Application Number
EP 93 12 0960

DOCUMENTS CONSIDERED TO BE RELEVANT

Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int.Cl.5)
A	IEEE INFOCOM '88, NEW ORLEANS, MARCH 27 - 31, 1988, no. 1988, 27 March 1988, INSTITUTE OF ELECTRICAL AND ELECTRONICS ENGINEERS, pages 9-18, XP000010853 AHMADI H ET AL: "A HIGH-PERFORMANCE SWITCH FABRIC FOR INTEGRATED CIRCUIT AND PACKET SWITCHING" * figure 2 * * page 11, column 2, line 41 - line 49 *	5	

TECHNICAL FIELDS SEARCHED (Int.Cl.5)

The present search report has been drawn up for all claims

Place of search	Date of completion of the search	Examiner
THE HAGUE	29 October 1996	Dhondt, E

CATEGORY OF CITED DOCUMENTS

X : particularly relevant if taken alone
Y : particularly relevant if combined with another document of the same category
A : technological background
O : non-written disclosure
P : intermediate document

T : theory or principle underlying the invention
E : earlier patent document, but published on, or after the filing date
D : document cited in the application
L : document cited for other reasons

& : member of the same patent family, corresponding document

EPO FORM 1503 03.82 (P04C01)